

# Threat models

# Basics of threat modelling

- Threat actors
- Adversaries
- Capabilities of adversaries
- Trust vs verifiability
- Clear articulation of all trust points

# Case study: Authentication and KYC

# Trust model of old-fashioned identity cards

- Presenter trusted?
- Verifier trusted?
- KYC based on identity documents?

# Trust model of old-fashioned identity cards

- Presenter trusted?
- Verifier trusted?
- KYC based on identity documents?
  - Possibilities of repurposing?
- Vacuous?

# Trust model of smart cards with chips

- Content trustworthy?
  - Under what conditions?
- Presenter?
- Verifier?
- Verifier machine?

# Trust model of Aadhaar Based Biometric Authentication

- No trust requirement on presenter?
- What about verifier?

# Trust model of Aadhaar Based Biometric Authentication

- No trust requirement on presenter?
- What about verifier?
  - Assume cannot control backend
  - False authorisation and/or accounting?
  - Store and replay?
- What if authentication outcome is routed through the verifier?

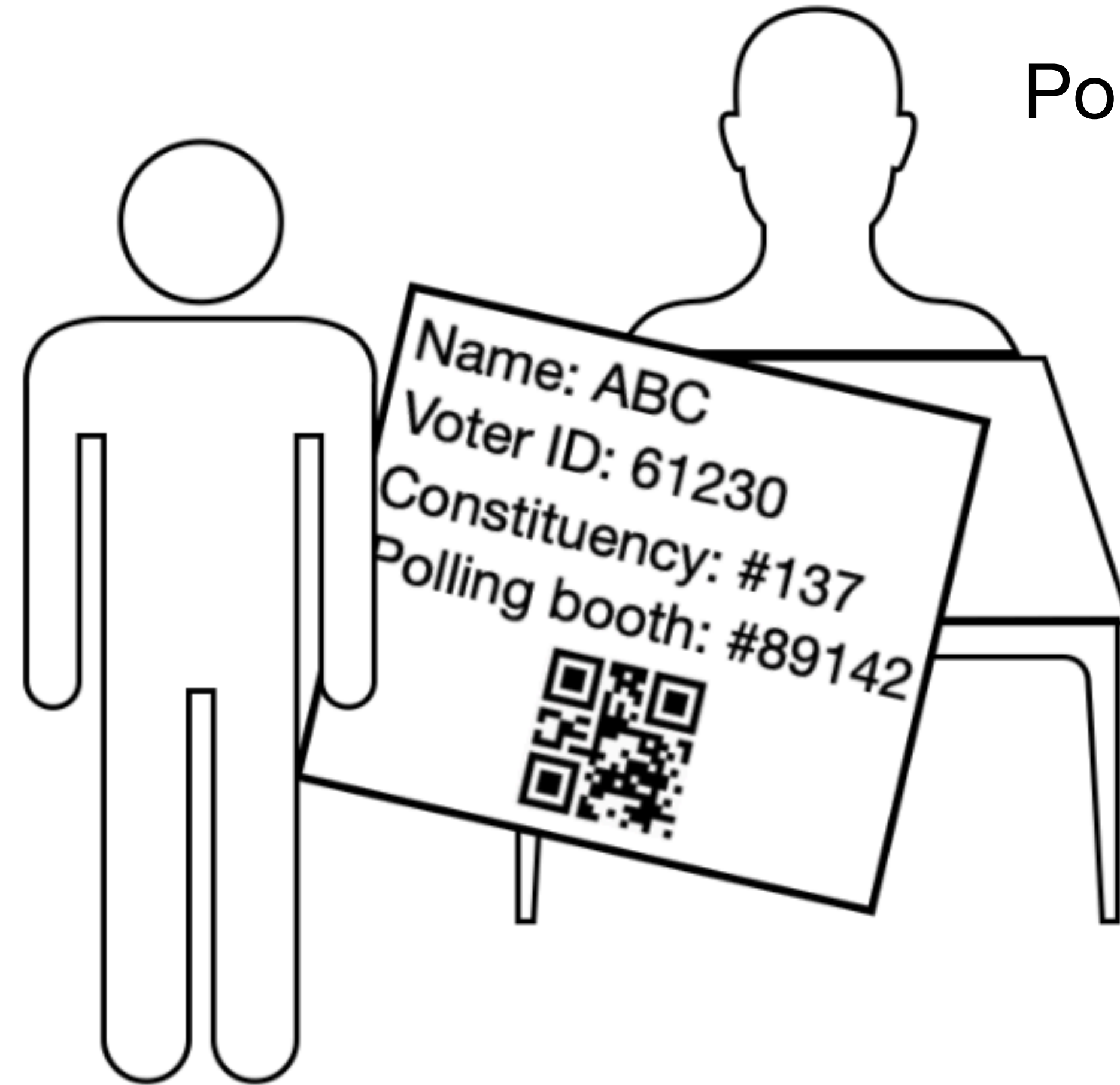


# Trust models of other authentication methods?

- Passwords
- Ssh authentication (Diffie-Helman key exchange)
- Kerberos authentication

# Case study: elections

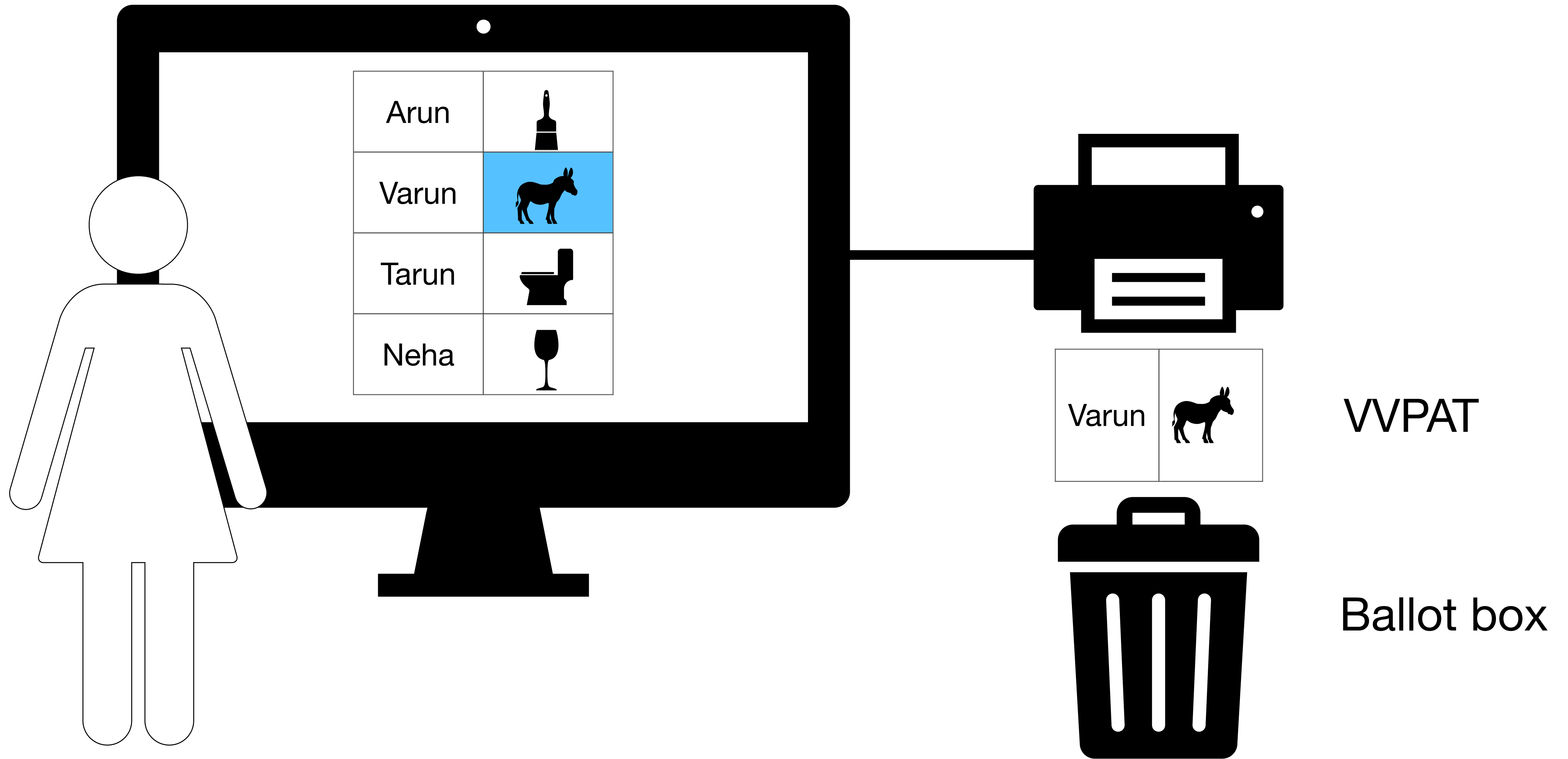
# Identity and eligibility verification



Polling officer

Also polling agents

# EVM + VVPAT



# **Security threat analysis**

# Threat model

- Adversary can corrupt and control
  - An arbitrary set of polling officials
  - An arbitrary set of voters
  - Voting equipment

# Voting requirements

- Correctness
  - Cast as intended
  - Recorded as cast
  - Counted as recorded
  - Only eligible voters and only 1 vote per eligible voters
  - Non-repudiation and dispute resolution
- Secrecy
  - Receipt free (voter should not be able to prove who she voted for)

# Formal definitions



# Threat model for verifiability

The adversary  $A_{\text{verifiability}}$  is a polynomially bounded adversary who may try to alter the outcome of the election. It

1. can corrupt the EA, the POs, the voting machines, or any other authorities;
2. can alter or delete cast votes during polling, during the collection and counting processes, or while publishing on public bulletin boards;
3. can introduce fake votes in the system, i.e., those not certified by polling officers;

# Verifiability

**Universal verifiability:** A voting system is universally verifiable if anyone in the public can verify using publicly posted data that

1. each vote is recorded-as-cast and counted-as-recorded
2. all recorded votes are cast by eligible voters
3. any eligible voter has cast at most one vote

**Individual verifiability:** A voting system is individually verifiable if any voter can obtain a sound proof that their vote is recorded-as-intended in the final tally

**Verifiability:** A voting system is verifiable if Universal verifiability and Individual verifiability hold in the presence of  $A_{\text{verifiability}}$

# Threat model for secrecy and coercion resistance

The adversary  $A_{\text{secrecy}}$  is a polynomially bounded adversary who may try to learn others' votes, or coerce them to vote in a certain way, or be a voter itself and try to prove to others how it voted. It

1. can observe all voter receipts, VVPRs (during counting) and public outputs posted on bulletin boards;
2. can participate as a bare-handed voter;
3. can interact with other voters before and after the voting process but not during;
4. can control POs and election authorities;
5. can corrupt voting machines to reveal votes or other secrets;
6. *cannot* observe secrets of to-be-used paper ballots between printing to their usage without leaving a trace of tampering;

# Secrecy and coercion resistance

**Individual vote secrecy:** A voting system protects individual vote secrecy if given a (possibly malicious) voter's receipt and publicly posted data, no information can be derived about how the voter voted. That is, a voter cannot prove to anybody how she voted.

**Community vote secrecy:** A voting system protects community vote secrecy if given voters' receipts and publicly posted data, an adversary cannot determine how voters assigned to a given polling booth voted.

**Secrecy preserving and coercion resistant:** A voting system is secrecy preserving and coercion resistant if the above two properties hold even in the presence of  $A_{\text{secrecy}}$

# **EVM+VVPAT system**

- Verifiable?
- Secrecy preserving and coercion resistant?

# EVM+VVPAT system

- Verifiable?
- Secrecy preserving and coercion resistant?
- **Software independent:** A protocol is software independent if an undetected change in the software cannot cause an undetectable change in the (election) outcome
- Software independence a necessary condition for both verifiability and secrecy?